



Meng-Meng Liu^{1,*,†}, Min-Qian Liu^{2,†} and Jin-Yu Yang^{2,†}

- ¹ School of Science, Minzu University of China, Beijing 100081, China
- ² NITFID, LPMC & KLMDASR, School of Statistics and Data Science, Nankai University, Tianjin 300071, China; mqliu@nankai.edu.cn (M.-Q.L.); jyyang@nankai.edu.cn (J.-Y.Y.)
- * Correspondence: liumengmeng@muc.edu.cn

[†] These authors contributed equally to this work.

Abstract: Computer experiments often involve both qualitative and quantitative factors, posing challenges for efficient experimental designs. Strongly coupled designs (SCDs) are proposed in this paper to balance flexibility in run size and stratification properties between qualitative and quantitative factor columns. The existence and construction of SCDs are investigated. When $s \ge 2$ is a prime or a prime power, the constructed SCDs of λs^3 runs can accommodate 2s - 1 qualitative factors and a substantial number of quantitative factors. Furthermore, a series of SCDs with s^u rows and $(u - 3)s^3$ columns of quantitative factors are constructed, where $u \ge 4$, with certain columns of quantitative factors achieving stratification in two or higher dimensions. The proposed SCDs have λs^4 rows, while SCDs have only λs^3 rows, offering more flexibility. Furthermore, in the designs constructed in this paper with fewer than 100 rows, in 11 out of 17 cases, SCDs have a larger number and higher levels of qualitative factors than DCDs.

Keywords: computer experiment; completely resolvable orthogonal array; qualitative and quantitative factor; regular design; space-filling design

MSC: 62K05; 62K99

1. Introduction

With the rapid development of computer science, an increasing number of scientists and engineers are using computer experiments to study complex physical systems. Spacefilling designs are feasible for these types of experiments as they spread the design points in the design region as uniformly as possible (Santner, Williams and Notz, 2003; Fang, Li and Sudjianto, 2006) [1,2]. Based on the effect sparsity principle (Wu and Hamada, 2021) [3], only a few factors are expected to be active, which makes it reasonable to focus on uniformity properties in low-dimensional projections. Latin hypercube designs (LHDs), which are among the most popular space-filling designs, satisfy one-dimensional uniformity (McKay, Beckman and Conover, 1979) [4]. Owen (1992) and Tang (1993) [5,6] proposed two methods to generate LHDs based on orthogonal arrays (OAs), which inherit *t*-dimensional projection properties when an OA of strength *t* is used. He and Tang (2013) [7] proposed a new class of arrays called strong orthogonal arrays (SOAs). To achieve better stratification properties, the strength of SOAs should be 3 or higher (He and Tang, 2014) [8], though this can lead to large run sizes. SOAs of strengths of 2+, introduced by He, Cheng and



Academic Editor: Yuhlong Lio Received: 28 November 2024 Revised: 25 December 2024 Accepted: 27 December 2024 Published: 28 December 2024

Citation: Liu, M.-M.; Liu, M.-Q.; Yang, J.-Y. Strongly Coupled Designs for Computer Experiments with Both Qualitative and Quantitative Factors. *Mathematics* **2025**, *13*, 75. https:// doi.org/10.3390/math13010075

Copyright: © 2024 by the authors. Licensee MDPI, Basel, Switzerland. This article is an open access article distributed under the terms and conditions of the Creative Commons Attribution (CC BY) license (https://creativecommons.org/ licenses/by/4.0/). Tang (2018) [9], are more economical than SOAs of strength 3 while retaining the same two-dimensional stratification. Zhou and Tang (2019) [10] proposed SOAs of strength 3–, which maintain the same three-dimensional stratification as SOAs of strength 3, as well as the same factor levels as SOAs of strength 2. Recent studies on the construction of designs achieving stratification in low dimensions also include those of Liu et al. (2024), and Wang, Lin and Liu (2024) [11,12].

In some cases, computer experiments involve both quantitative and qualitative factors. Qian, Wu and Wu (2008) [13] proposed a Gaussian process model framework to address the integration of qualitative and quantitative factors in computer modeling. This framework achieves integration by constructing and validating correlation functions that incorporate both types of factors. Han et al. (2009) [14] proposed a Bayesian methodology for predicting computer experiments with both qualitative and quantitative inputs, modeling outputs for different qualitative levels as having similar functional behavior in the quantitative inputs. Huang et al. (2016) [15] introduced SLHDs with good uniformity for each slice and an adaptive analysis strategy to improve prediction precision and assess similarities among qualitative variable levels in computer experiments.

Qian (2012) [16] suggested sliced LHDs (SLHDs) for this type of experiment. Recently, Kumar et al. (2024) [17] investigated the construction methods for sliced orthogonal LHDs with unequal batch sizes. Guo et al. (2023) [18] proposed methods for constructing orthogonal and nearly orthogonal general SLHDs; these designs ensure orthogonality not only in the overall design but also within each layer before and after collapsing. Wang, Wang and Xue (2023) [19] proposed a sequential optimal LHD method to improve the spacefilling properties of two-layer computer simulators. In SLHDs, stratification is achieved between all qualitative and quantitative factors; however, this can lead to a significant increase in the run sizes as the number and levels of qualitative factors expand. To provide more economical run sizes, Deng, Hung and Lin (2015) [20] proposed marginally coupled designs (MCDs). Subsequently, a large number of researchers have contributed to the study of MCDs with better properties, such as orthogonality or low-dimensional stratification, including He, Lin and Sun (2017); He et al. (2017); He, Lin and Sun (2019); and Zhou, Yang and Liu (2021) [21–24]. To study the interaction of the qualitative and quantitative factors, Yang et al. (2023) [25] proposed doubly coupled designs (DCDs) and studied their constructions. DCDs employ better stratification properties between qualitative and quantitative factors compared to MCDs, though this comes at the cost of reducing the number of qualitative factors. Zhou, Huang and Li (2024) [26] investigated the construction of group DCDs.

In this paper, a new type of design called strongly coupled designs (SCDs) is proposed for computer experiments with both quantitative and qualitative factors. Similar to DCDs, SCDs achieve stratification between any two qualitative factors and all quantitative factors to a certain extent. Compared to MCDs, SCDs are more feasible for studying interactions between any two qualitative factors and all quantitative factors. The stratification requirements among qualitative factors are relaxed which makes the run sizes of SCDs more flexible than those of DCDs. For example, when the levels of the qualitative factors are s^2 , DCDs divide the design corresponding to the quantitative factors into s^4 small LHDs, with each corresponding to a specific level combination of two qualitative factors, while SCDs divide it into s^3 small LHDs.

The characteristics and construction methods of SCDs are also studied. The necessary and sufficient conditions for the existence of SCDs are provided from three different perspectives, and SCDs with λs^3 runs, where $s \ge 2$ is a prime or a prime power, and λ is a positive integer, are constructed. For a given run size, the constructed SCDs can accommodate a greater number of qualitative factors than DCDs in many cases. Additionally,

the construction of SCDs with desirable space-filling properties for quantitative factors is investigated. A series of designs with s^u runs and $(u - 3)s^3$ columns of quantitative factors, where $u \ge 4$, are constructed. In these designs, certain quantitative factor columns achieve stratification in two or more dimensions.

The remainder of this paper is organized as follows. In Section 2, some basic concepts are introduced, SCDs are defined, and their existence is studied. In Section 3, the construction of SCDs is explored, including those with desirable space-filling properties for quantitative factors. Examples are provided to illustrate the construction process and the properties of the resulting designs. The paper is concluded with a discussion in Section 4. Some theorem proofs are deferred to the Appendix A.

2. Definitions and Characterizations

In this section, key concepts relevant to our work are reviewed, and the definition of SCDs is proposed. Additionally, three characterizations of SCDs are presented, providing heuristic guidance for the construction methods discussed in the next section.

An $n \times m$ array, where the *j*-th column has s_j levels, $\{0, 1, \ldots, s_j - 1\}$, is called an OA of strength *t*, denoted by OA $(n, m, s_1 \times \cdots \times s_m, t)$, if every $n \times t$ subarray contains each *t*-tuple with equal frequency. When $s_1 = \cdots = s_m = s$, the array is symmetric and is denoted as OA(n, m, s, t) for simplicity. An OA(n, m, s, t) is said to be completely resolvable, denoted by CROA(n, m, s, t), if its rows can be partitioned into n/s subarrays, each of which is an OA(s, m, s, 1).

An $r \times e$ array D(r, e, s) is called a difference scheme based on GF(s) if it satisfies the following property: for all i and j such that $1 \leq i, j \leq e$ and $i \neq j$, the vector difference between the *i*-th and *j*-th columns contains each element of GF(s) with equal frequency. Difference schemes are fundamental in the construction of CROAs.

An SOA of strength *t*, denoted by SOA(n, m, s^t, t), is an $n \times m$ matrix with entries from {0,1,..., $s^t - 1$ }. For any subarray consisting of *g* columns ($1 \le g \le t$), it can be collapsed into an OA($n, g, s^{u_1} \times \cdots \times s^{u_g}, g$), where u_1, \ldots, u_g are positive integers satisfying $u_1 + \cdots + u_g = t$. The s^t levels of a factor are collapsed into s^{u_j} levels using the mapping $\lfloor a/s^{t-u_j} \rfloor$, where $a = 0, 1, \ldots, s^t - 1$ and $1 \le j \le g$. Here, $\lfloor b \rfloor$ denotes the greatest integer less than or equal to *b*. An SOA of strength 2+, denoted by SOA($n, m, s^2, 2+$), is an SOA($n, m, s^2, 2$) with the additional property that any two distinct columns can be collapsed into both an OA($n, 2, s^2 \times s, 2$) and an OA($n, 2, s \times s^2, 2$). Similarly, an SOA of strength 3-, denoted by SOA($n, m, s^2, 3-$), is an SOA($n, m, s^2, 2+$) with the property that any three distinct columns can be collapsed into an OA(n, 3, s, 3).

For two matrices $E = (e_{ij})_{m \times n}$ and $F = (f_{ij})_{u \times v}$ with entries from GF(s), their Kronecker sum \oplus is defined as

$$E \oplus F = \begin{pmatrix} e_{11} + F & \cdots & e_{1n} + F \\ \vdots & \ddots & \vdots \\ e_{m1} + F & \cdots & e_{mn} + F \end{pmatrix},$$

where $\dot{+}$ denotes addition in GF(s). 1_s is used to denote a vector of length s consisting entirely of ones. For any vector $w = (w_1, \ldots, w_m)^T$, $w \otimes 1_s$ is defined as $(w_1, \ldots, w_1, w_2, \ldots, w_m)^T$ and $1_s \otimes w$ is defined as $(w_1, \ldots, w_m, \ldots, w_1, \ldots, w_m)^T$.

Let $D = (D_1, D_2)$ be a design with q qualitative factors and p quantitative factors, where D_1 and D_2 are sub-designs for qualitative and quantitative factors, respectively. A design D is called an MCD if D_2 is an LHD and the rows in D_2 corresponding to each level of any factor in D_1 form a small LHD. A DCD is an MCD with the additional property that the rows in D_2 corresponding to each level combination of any two factors in D_1 form an LHD.

In this paper, we introduce a new class of designs for computer experiments with both quantitative and qualitative factors. Suppose the s^2 levels of any qualitative factor are divided into *s* groups based on a specified grouping method, with each group labeled by a group number: 0, 1, ..., s - 1. The definition of SCDs is presented below.

Definition 1. An n-run design $D = (D_1, D_2)$, with $q s^2$ -level qualitative factors and p quantitative factors, is called a strongly coupled design, denoted by $SCD(n, (s^2)^q, p)$, if it satisfies the following conditions:

- (i) D_2 is an LHD;
- (ii) The rows in D_2 corresponding to each group number of any column in D_1 form an LHD;
- (iii) The rows in D_2 corresponding to each combination of group number and level from any two columns in D_1 form an LHD.

For $D_1 = (z_1, ..., z_q)$, to simplify the following discussion, the s^2 levels of qualitative factors are encoded as $0, 1, ..., s^2 - 1$. The group number is defined as $\bar{z}_i = \lfloor \frac{z_i}{s} \rfloor$, which is obtained by collapsing the s^2 levels into s levels, with levels sharing the same group number classified into the same group. In this paper, D_1 is set as an SOA($n, q, s^2, 2+$). For $D_2 = (d_1, ..., d_p)$, define

$$\overline{D}_2 = \lfloor \frac{D_2}{s} \rfloor = (\overline{d}_1, \dots, \overline{d}_p) \text{ and } \widetilde{D}_2 = \lfloor \frac{\overline{D}_2}{s^2} \rfloor = (\widetilde{d}_1, \dots, \widetilde{d}_p),$$

where |a| represents the largest integer not exceeding *a*.

Next, we present the necessary and sufficient conditions for the existence of SCDs. Theorem 1 specifies the conditions that must be satisfied between columns from D_1 and D_2 .

Theorem 1. Suppose $D_1 = (z_1, \ldots, z_q)$ is an SOA $(n, q, s^2, 2+)$, and $D_2 = (d_1, \ldots, d_p)$ is an LHD(n, p). The design $D = (D_1, D_2)$ is an SCD $(n, (s^2)^q, p)$ if and only if:

- (i) (\bar{z}_i, \bar{d}_k) is an OA(n, 2, s(n/s), 2) for any $1 \le i \le q, 1 \le k \le p$;
- (ii) $(z_i, \bar{z}_j, \tilde{d}_k)$ is an OA $(n, 3, s^2 s(n/s^3), 3)$ for any $1 \le i \ne j \le q, 1 \le k \le p$.

Assume that D_1 is generated by the formula sA + B. Theorem 2 provides the necessary and sufficient conditions that the columns of A and B must satisfy for the existence of SCDs.

Theorem 2. Suppose that $D_1 = (z_1, ..., z_q)$ is an SOA $(n, q, s^2, 2+)$ generated by $D_1 = sA + B$, where $A = (\overline{z}_1, ..., \overline{z}_q)$ and $B = (b_1, ..., b_q)$. The design $D = (D_1, D_2)$ is an SCD $(n, (s^2)^q, p)$ if and only if:

- (i) A can be divided as $(A^{(1)T}, \ldots, A^{(n/s^3)T})^T$, where each $A^{(l)}$ is a CROA $(s^3, q, s, 2)$, for $1 \leq l \leq n/s^3$;
- (ii) For any two columns $\bar{z}_i^{(l)}$ and $\bar{z}_j^{(l)}$ in $A^{(l)}$, and the corresponding rows $b_i^{(l)}$ in b_i , the three tuples $(\bar{z}_i^{(l)}, \bar{z}_j^{(l)}, b_i^{(l)})$ form an OA $(n, 3, s^3, 3)$, for any $1 \le i \ne j \le q, 1 \le l \le n/s^3$.

Assume that $D_1 = sA + B$ and $\overline{D}_2 = s^2C + sE + F$. Theorem 3 provides the necessary and sufficient conditions that the columns of *A*, *B*, *C*, *E*, and *F* must satisfy.

Theorem 3. Suppose D_1 is an SOA $(n, q, s^2, 2+)$ and D_2 is an LHD(n, p), where D_1 and \overline{D}_2 are generated by $D_1 = sA + B$ and $\overline{D}_2 = s^2C + sE + F$, respectively. The design $D = (D_1, D_2)$ is an SCD $(n, (s^2)^q, p)$ if and only if there exist five arrays: $A = (\overline{z}_1, \dots, \overline{z}_q) = OA(n, q, s, 2)$, $B = (b_1, \dots, b_q) = OA(n, q, s, 1)$, $C = (\tilde{d}_1, \dots, \tilde{d}_p) = OA(n, p, n/s^3, 1)$, $E = (e_1, \dots, e_p) =$

OA(n, p, s, 1), and $F = (f_1, \dots, f_q) = OA(n, q, s, 1)$, such that for any $1 \le i \ne j \le q$ and $1 \le k \le p$, both $(\overline{z}_i, f_k, e_k, \widetilde{d}_k)$ and $(\overline{z}_i, b_i, \overline{z}_j, \widetilde{d}_k)$ are OA $(n, 4, s^3(n/s^3), 4)$.

3. Construction Methods

In this section, methods for constructing SCDs, as well as the ones with desirable spacefilling properties for quantitative factors, are proposed. We focus on cases where $s \ge 2$ is a prime or a prime power. The SCDs constructed in Section 3.1 can accommodate 2s - 1qualitative factors and a large number of quantitative factors. For the SCDs constructed in Section 3.2, when the run size $n = s^u$, the number of columns for quantitative factors is $(u - 3)s^3$, and certain columns of quantitative factors achieve stratification in two or more higher dimensions.

3.1. Construction of SCDs

SCDs with run sizes $n = \lambda s^3$ can be obtained using Algorithm 1, where $s \ge 2$ is a prime or a prime power, and λ is a positive integer.

Algorithm 1: Construction of SCDs

Step 1.	Let $H = (h_1,, h_s)$ be a $D(s, s, s)$. Define $O = (0,, s - 1)^T \oplus H = (o_1,, o_s)$,
	an OA obtained by developing the difference schemes H .
Step 2.	Define $a_1 = h_1 \oplus o_1$, and $A_i = (h_1, h_i) \oplus o_i$ for $i = 2, \dots, s$. Let
	$A^{(1)} = \cdots = A^{(\lambda)} = (a_1, A_2, \dots, A_s), \text{ and } A = (A^{(1)^T}, \dots, A^{(\lambda)^T})^T.$
Step 3.	For $1 \leq j_{i,m}$, $k_{i,m} \leq s$, where $1 \leq i \leq s$, $1 \leq m \leq 2$, $j_{i,m} \neq k_{i,m}$ and $j_{i,m} \neq 1$, define:
	$b_1 = h_{j_{1,1}} \oplus o_{k_{1,1}}$ with $j_{1,1} \neq 1$; $b_{i1} = h_{j_{i,1}} \oplus o_{k_{i,1}}$ with $k_{i,1} \neq 1$ and $j_{i,1} \neq i$, or $k_{i,1} = 1$
	and $j_{i,1} = i$; $b_{i2} = h_{j_{i,2}} \oplus o_{k_{i,2}}$ with $j_{i,2} = i$. Let $B_i = (b_{i1}, b_{i2})$, and
	$B^{(1)} = \cdots, B^{(\lambda)} = (b_1, B_2, \dots, B_s).$ Then obtain $B = (B^{(1)T}, \dots, B^{(\lambda)T})^T.$
Step 4.	Define $c_i = \omega_i \otimes 1_{s^3}$, where ω_i is a permutation of $(0,, \lambda - 1)$ for $1 \le i \le \lambda!$. For
	$1 \leq l \leq \lambda$ and $1 \leq k \leq s^2$!, define $e_k^{(l)} = ((1_s \otimes \beta_{1,k}^{(l)})^T, \dots, (1_s \otimes \beta_{s,k}^{(l)})^T)^T$, where
	$(\beta_{1,k}^{(l)T},\ldots,\beta_{s,k}^{(l)T})^T$ is a permutation of $(0,\ldots,s^2-1)^T$. For $1 \le \delta \le (s^2!)^{\lambda}$, let
	$e_{\delta} = (e_{k_1}^{(1)^T}, \dots, e_{k_{\lambda}}^{(\lambda)^T})^T$, where $1 \leq k_l \leq s^2$!. Then, obtain $E = (e_1, \dots, e_{(s^2!)^{\lambda}})$.
Step 5.	Let $D_1 = sA + B$ and $\overline{D}_2 = (s^2c_1 + E, \dots, s^2c_{\lambda!} + E)$. For each column in \overline{D}_2 , replace
	the <i>s</i> positions of level <i>i</i> by all possible random permutation of $\{is, is + 1, \ldots, interpretext \}$
	$is + (s - 1)$, $i = 0, 1,, n/s - 1$, columns juxtapose these matrices to obtain D_2 ,
	then obtain $D = (D_1, D_2)$.

To make it easier for readers to understand the algorithm, we provide the flow chart in Figure 1 to explain the algorithm.

The properties of the constructed designs are summarized in the following theorem.

Theorem 4. The design D constructed in Algorithm 1 is an SCD $(n, (s^2)^q, p)$ with $n = \lambda s^3$, q = 2s - 1 and $p = \lambda ! (s^2!)^{\lambda} (s!)^{\lambda s^2}$, where λ is a positive integer.

An example is now given to illustrate the construction procedure and the properties of the constructed designs.



Figure 1. Flow chart of Algorithm 1.

Example 1. For $s = \lambda = 2$, let

$$H = (h_1, h_2) = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix},$$

then,

$$O = (0,1)^T \oplus H = \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{pmatrix}^T = (o_1, o_2).$$

Define $a_1 = h_1 \oplus o_1 = (0, 0, 1, 1, 0, 0, 1, 1)^T$,

$$A_2 = (h_1, h_2) \oplus o_2 = \begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 1 \end{pmatrix}^T,$$

then,

$$A = \begin{pmatrix} A^{(1)} \\ A^{(2)} \end{pmatrix} = \begin{pmatrix} a_1 & A_2 \\ a_1 & A_2 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 1 \end{pmatrix}^T.$$

Define $b_1 = b_{21} = b_{22} = h_2 \oplus o_1 = (0, 0, 1, 1, 1, 1, 0, 0)^T$, then,

$$B = \begin{pmatrix} B^{(1)} \\ B^{(2)} \end{pmatrix} = \begin{pmatrix} b_1 & B_2 \\ b_1 & B_2 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \end{pmatrix}^T$$

As $\lambda = 2$, let $w_1 = (0, 1)$ and $w_2 = (1, 0)$. Correspondingly, obtain

$$c_2 = w_2 \otimes 1_8 = (1, 1, 1, 1, 1, 1, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0)^T$$

For $1 \leq k \leq 24$, both $(\beta_{1,k}^{(1)^T}, \beta_{2,k}^{(1)^T})^T$ and $(\beta_{1,k}^{(2)^T}, \beta_{2,k}^{(2)^T})^T$ are a permutation of $(0, 1, 2, 3)^T$. Let

$$e_{\delta} = \begin{pmatrix} e_{k_{1}}^{(1)} \\ e_{k_{2}}^{(2)} \end{pmatrix} = \begin{pmatrix} (1_{2} \otimes \beta_{1,k_{1}}^{(1)})^{T}, & (1_{2} \otimes \beta_{2,k_{1}}^{(1)})^{T} \\ (1_{2} \otimes \beta_{1,k_{2}}^{(2)})^{T}, & (1_{2} \otimes \beta_{2,k_{2}}^{(2)})^{T} \end{pmatrix}^{T},$$

for $1 \leq \delta \leq 24^2$, $1 \leq k_1, k_2 \leq 24$. Define $E = (e_1, \ldots, e_{24^2})$, then, obtain $\overline{D}_2 = (4c_1 + E, 4c_2 + E)$. The first 24 columns of \overline{D}_2 are listed in Table 1. The two positions of level $0, 1, \ldots, 7$ are replaced by all possible random permutations of $\{0, 1\}, \{2, 3\}, \ldots, \{14, 15\}$, respectively. The columns of these matrices are then juxtaposed to form D_2 .

The resulting design $D = (D_1, D_2)$ is an $SCD(16, 4^3, 24^22^9)$. When collapsing D_1 into 2 levels, the rows in D_2 corresponding to each level of any factor in D_1 form an LHD(8, 24²2⁹). When collapsing any two columns of D_1 into 4×2 or 2×4 levels, the rows in D_2 corresponding to each level combination of any two factors in D_1 form an LHD(2, 24²2⁹). Furthermore, the first two columns of D_2 achieve stratification on 2×2 grids.

Table 1. The SCD generated in Example 1.

	D_1				The First 24 Columns of $ar{D}_2$																					
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	2	2	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
3	3	3	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
3	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
1	1	3	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
1	3	1	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
2	2	0	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	2	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
0	0	0	4	5	5	4	6	6	5	4	6	6	4	5	6	6	4	5	5	4	7	7	7	7	7	7
0	2	2	5	4	6	6	4	5	6	6	4	5	5	4	7	7	7	7	7	7	4	5	5	4	6	6
3	3	3	4	5	5	4	6	6	5	4	6	6	4	5	6	6	4	5	5	4	7	7	7	7	7	7
3	1	1	5	4	6	6	4	5	6	6	4	5	5	4	7	7	7	7	7	7	4	5	5	4	6	6
1	1	3	6	6	4	5	5	4	7	7	7	7	7	7	4	5	5	4	6	6	5	4	6	6	4	5
1	3	1	7	7	7	7	7	7	4	5	5	4	6	6	5	4	6	6	4	5	6	6	4	5	5	4
2	2	0	6	6	4	5	5	4	7	7	7	7	7	7	4	5	5	4	6	6	5	4	6	6	4	5
2	0	2	7	7	7	7	7	7	4	5	5	4	6	6	5	4	6	6	4	5	6	6	4	5	5	4

Table 2 lists various SCDs generated for different positive integer values of λ , demonstrating the flexibility of SCDs in accommodating different run sizes. Each row in the table shows the design's ability to accommodate qualitative and quantitative factors. As the ex-

periment size increases, the number of quantitative factors that the design can accommodate increases significantly.

s	п	D_1	D_2
2	8λ	$SOA(8\lambda, 3, 2^2, 2+)$	LHD $(8\lambda, \lambda! (4!)^{\lambda} 2^{4\lambda})$
3	27λ	$SOA(27\lambda, 5, 2^2, 2+)$	LHD(27 λ , λ !(9!) $^{\lambda}$ (3!) $^{9\lambda}$)
4	64λ	$SOA(64\lambda, 7, 2^2, 2+)$	LHD(64 λ , λ !(16!) $^{\lambda}$ (4!) ^{16λ})
5	125λ	$SOA(125\lambda, 9, 2^2, 2+)$	LHD($125\lambda, \lambda!(25!)^{\lambda}(5!)^{25\lambda}$)

Table 2. Some generated SCDs where λ is a positive integer.

Table 3 presents the comparison of SCDs, DCDs, and MCDs with the same run sizes. All three kinds of designs can accommodate a large number of qualitative factors, with MCDs supporting significantly more qualitative and quantitative factors than SCDs and DCDs. However, SCDs and DCDs outperform MCDs in terms of the stratification property between qualitative and quantitative factors. For SCDs and DCDs, in 11 out of 17 cases, SCDs have a larger number and higher levels of qualitative factors than DCDs. In the remaining cases, SCDs accommodate only one fewer qualitative factor than DCDs.

Table 3. Comparison of SCDs, DCDs, and MCDs.

п	SCD	DCD	MCD
8	SCD(8,4 ³ ,4!2 ⁴)	DCD(8,2 ² ,2 ⁶ 2!)	MCD(8,2 ⁴ ,4!2 ⁴)
16	$SCD(16,4^3,2!(4!)^22^8)$	$DCD(16,4^4,(4!)^5)$	MCD(16,4 ⁴ ,4!24 ⁴)
24	$SCD(24,4^3,3!(4!)^32^{12})$	DCD(24,2 ² ,2 ¹⁸ 6!)	MCD(24,2 ¹² ,12!2 ¹²)
27	SCD(27,9 ⁵ ,9!(3!) ⁹)	DCD(27,3 ³ ,6 ¹² 3!)	MCD(27,3 ⁹ ,9!6 ⁹)
32	SCD(32,4 ³ ,4!(4!) ⁴ 2 ¹⁶)	$DCD(32,4^4,(4!)^{10}2!)$	MCD(32,4 ⁸ ,8!24 ⁸)
40	$SCD(40,4^3,5!(4!)^52^{20})$	DCD(40,2 ² ,2 ³⁰ 10!)	MCD(40,2 ²⁰ ,20!2 ²⁰)
48	SCD(48,4 ³ ,6!(4!) ⁶ 2 ²⁴)	DCD(48,4 ⁴ ,(4!) ¹⁵ 3!)	MCD(48,4 ¹² ,12!24 ¹²)
54	$SCD(54,9^5,2!(9!)^2(3!)^{18})$	DCD(54,3 ³ ,6 ²⁴ 6!)	MCD(54,3 ¹⁸ ,18!6 ¹⁸)
56	$SCD(56,4^3,7!(4!)^72^{28})$	DCD(56,2 ² ,2 ⁴² 14!)	MCD(56,2 ²⁸ ,28!2 ²⁸)
64	SCD(64,4 ³ ,8!(4!) ⁸ 2 ³²)	DCD(64,4 ⁴ ,(4!) ²⁰ 4!)	MCD(64,4 ¹⁶ ,16!24 ¹⁰)
72	$SCD(72,4^3,9!(4!)^92^{36})$	DCD(72,3 ³ ,6 ³² 8!)	MCD(72,3 ²⁴ ,24!6 ²⁴)
80	SCD(80,4 ³ ,10!(4!) ¹⁰ 2 ⁴⁰)	$DCD(80,4^4,(4!)^{25}5!)$	MCD(80,4 ⁸ ,20!24 ²⁰)
81	$SCD(81,9^5,3!(9!)^3(3!)^{27})$	DCD(81,3 ³ ,6 ³⁶ 9!)	MCD(81,3 ²⁷ ,27!6 ²⁷)
88	$SCD(88,4^3,11!(4!)^{11}2^{44})$	DCD(88,2 ² ,2 ⁶⁶ 22!)	MCD(88,2 ⁴⁴ ,44!2 ⁴⁴)
96	SCD(96,4 ³ ,12!(4!) ¹² 2 ⁴⁸)	DCD(96,4 ⁴ ,(4!) ³⁰ 6!)	MCD(96,4 ¹⁶ ,24!24 ²⁴)

3.2. Construction of SCDs with Desirable Space-Filling Properties

This section focuses on constructing SCDs in which D_2 has desirable space-filling properties. A series of designs with s^u rows and $(u - 3)s^3$ columns of quantitative factors are constructed for $u \ge 4$. These designs ensure that certain columns of quantitative factors achieve stratification in two or higher dimensions.

Let ξ_1, \ldots, ξ_u represent u independent columns with entries from GF(s), each of length $n = s^u$. A regular saturated OA can be constructed by first listing these independent columns and then adding all possible interactions. This design is denoted as S = $\{u_1\xi_1 + \cdots + u_u\xi_u \mid u_i \in GF(s), (u_1, \ldots, u_u) \neq (0, \ldots, 0), \text{ and the first nonzero entry is 1}\}$. Subsequently, select columns from S to construct SCDs in which D_2 possesses desirable space-filling properties. For when $s \ge 3$ is a prime or a prime power, we propose our construction method in Algorithm 2.

Algorithm 2: Construction of SCDs with desirable space-filling properties

Step 1. Define $A_1 = \{\xi_1 + (s-1)\xi_2 + w\xi_3 \mid 1 \le w \le s-2\}$ and $A_2 = \{\xi_1 + w\xi_2 + (s-1)\xi_3 \mid 1 \le w \le s-2\}$. Construct

$$A = (\xi_1 + (s-1)\xi_2, \xi_2 + (s-1)\xi_3, A_1, A_2, \xi_1 + (s-1)\xi_2 + (s-1)\xi_3).$$

- Step 2. Repeat ξ_2 and $\xi_3 s 2$ times to form $B_1 = (\xi_2, \dots, \xi_2)$ and $B_2 = (\xi_3, \dots, \xi_3)$, respectively. Let $B = (\xi_1, \xi_2, B_1, B_2, \xi_1)$. Construct $D_1 = sA + B$.
- Step 3. Define $E = (\xi_1 + (s 1)\xi_2, ..., \xi_1 + (s 1)\xi_2)$ and $F = (\xi_1, ..., \xi_1)$, where $\xi_1 + (s 1)\xi_2$ and ξ_1 are each repeated $(u 3)s^3$ times.

Step 4. Let $R_v = \{\xi_1 + u_2\xi_2 + u_3\xi_3 + u_{v+3}\xi_{v+3}|u_2, u_3 \in GF(s), u_{v+3} \in GF(s) \setminus \{0\}\} \cup \{\xi_2 + u_3\xi_3 + u_{v+3}\xi_{v+3}|u_3 \in GF(s), u_{v+3} \in GF(s) \setminus \{0\}\} \cup \{\xi_3 + u_{v+3}\xi_{v+3}|u_{v+3} \in GF(s) \setminus \{0\}\} \cup \{\xi_{v+3}\} = (r_{v,1}, \dots, r_{v,s^3})$, where $1 \le v \le u - 3$. Define

$$T = \begin{pmatrix} s^{u-4} & 1 & \cdots & s^{u-6} & s^{u-5} \\ s^{u-5} & s^{u-4} & \cdots & s^{u-7} & s^{u-6} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ s & s^2 & \cdots & s^{u-4} & 1 \\ 1 & s & \cdots & s^{u-5} & s^{u-4} \end{pmatrix} = (t_1, \dots, t_{u-3}).$$

Let $C = (C_1, ..., C_{s^3})$, where

$$C_f = (r_{1,f}, \dots, r_{u-3,f})T, \quad 1 \leq f \leq s^3.$$

Step 5. Construct $\overline{D}_2 = s^2 C + sE + F$. For each column of \overline{D}_2 , replace the *s* positions of

level *i* with a random permutation of $\{is, is + 1, ..., is + (s - 1)\}$ to obtain D_2 , where $i = 0, 1, ..., s^2 - 1$. Finally, let $D = (D_1, D_2)$.

A flow chart is provided in Figure 2 to explain the algorithm.

For s = 2, SCDs with desirable space-filling properties can be constructed by modifying Algorithm 2 as follows: set $A = (\xi_1, \xi_1 + \xi_2)$, $B = (\xi_1 + \xi_2 + \xi_3, \xi_1 + \xi_3)$, $E = (\xi_2, ..., \xi_2)$, and $F = (\xi_1 + \xi_3, ..., \xi_1 + \xi_3)$, where ξ_2 and $\xi_1 + \xi_3$ are repeated $(u - 3)s^3$ times. The properties of the resulting designs for $s \ge 2$, where s is a prime or a prime power, are summarized in Theorem 5.

Theorem 5. For $s \ge 2$, where *s* is a prime or a prime power, and $u \le 4$, the design *D* constructed above is an SCD $(s^u, (s^2)^q, (u-3)s^3)$ with the following properties:

(i)

 $q = \begin{cases} 2, & \text{when } s = 2, \\ 2s - 1, & \text{when } s \ge 3 \text{ is a prime or a prime power;} \end{cases}$

- (ii) For any *m* columns $\tilde{d}_{i_1}, \ldots, \tilde{d}_{i_m}$ with $\lfloor (i_1 1)/(u 3) \rfloor = \cdots = \lfloor (i_m 1)/(u 3) \rfloor$, where $m \leq u - 3$, the columns $\tilde{d}_{i_1}, \ldots, \tilde{d}_{i_m}$ achieve stratification on $s \times \cdots \times s$ grids;
- (iii) When $\lfloor (i-1)/(u-3) \rfloor \neq \lfloor (i'-1)/(u-3) \rfloor$, the columns \tilde{d}_i and $\tilde{d}_{i'}$ achieve stratification on $s^{u-3} \times s$ and $s \times s^{u-3}$ grids.



Figure 2. Flow chart of Algorithm 2.

An example is now provided to illustrate the construction procedure and the properties of the resulting designs.

Example 2. For s = 3 and u = 4, let $\xi_1, \xi_2, \xi_3, \xi_4$ be four independent columns of length 3^4 . Define $A_1 = \xi_1 + 2\xi_2 + \xi_3$ and $A_2 = \xi_1 + \xi_2 + 2\xi_3$. Construct

$$A = (\xi_1, \xi_2, \xi_3) \begin{pmatrix} 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 2 & 1 & 2 \\ 2 & 2 & 1 & 2 & 2 \end{pmatrix}$$

Let $B_1 = \xi_2$ and $B_2 = \xi_3$. Then,

$$B = (\xi_1, \xi_2, \xi_3) \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}.$$

Construct $D_1 = sA + B$; the resulting D_1 is presented in Table 4.

Next, define $E = (\xi_1 + 2\xi_2, ..., \xi_1 + 2\xi_2)$ *and* $F = (\xi_1, ..., \xi_1)$ *, where both* $\xi_1 + 2\xi_2$ *and* ξ_1 *are repeated nine times. Let* $R_1 = (\xi_1, \xi_2, \xi_3, \xi_4)$

/1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	0	0	0	0	0	0	0	0/	
0	1	2	0	1	2	0	1	2	0	1	2	0	1	2	0	1	2	1	1	1	1	1	1	0	0	0	
0	0	0	1	1	1	2	2	2	0	0	0	1	1	1	2	2	2	0	1	2	0	1	2	1	1	0	1.
$\backslash 1$	1	1	1	1	1	1	1	1	2	2	2	2	2	2	2	2	2	1	1	1	2	2	2	1	2	1/	

With T = 1, we have $C = R_1$. Construct $\overline{D}_2 = s^2C + sE + F$, which is also presented in Table 4. Finally, obtain D_2 by replacing the three positions of each level $0, 1, \ldots, 26$ with random permutations of $\{0, 1, 2\}, \{3, 4, 5\}, \ldots, \{78, 79, 80\}$, respectively.

 Table 4. The generated SCD in Example 2.

50066633300066663330006666333444111777744441117777444411177778888555222288855522228885555222288885555222288855522228885552222888555222288855522222888555222228885552222288855522222888555222228885552222288855522222888555222228885552222288855522222888555222228885552282852828528282282	
000666633344411177788855522220006666333444111777888555222200066663334441111777888555522220006666333444111177788855552222000666633344411117778885555222200066663334444111177788855552222200066666333444411117778885555222220006666633344441111777888555522222000066666333344441111777888555522222000666666666666	
003336666777111144455558882222333666600011114444777888822225556666000333344447771111222555556666000333344447777111122255555555555666600033334444777711112222555555555555555555555555	$\frac{D_1}{c}$
$0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 3 \\ 3 \\ 1 \\ 1 \\ 1 \\ 8 \\ 8 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 3 \\ 3 \\ 1 \\ 1 \\ 1 \\ 8 \\ 8 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 7 \\ 7 \\ 7 \\ 5 \\ 5 \\ 6 \\ 6 \\ 6 \\ 4 \\ 4 \\ 4 \\ 2 \\ 2 \\ 2 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0$	0
$0\\ 0\\ 0\\ 6\\ 6\\ 6\\ 3\\ 3\\ 3\\ 0\\ 0\\ 0\\ 0\\ 3\\ 3\\ 0\\ 0\\ 0\\ 0\\ 6\\ 6\\ 6\\ 4\\ 4\\ 4\\ 4\\ 1\\ 1\\ 1\\ 7\\ 7\\ 7\\ 1\\ 1\\ 1\\ 7\\ 7\\ 7\\ 4\\ 4\\ 4\\ 4\\ 7\\ 7\\ 7\\ 4\\ 4\\ 4\\ 1\\ 1\\ 1\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 2\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 8\\ 5\\ 5\\ 5\\ 2\\ 2\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 2\\ 2\\ 2\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\ 8\\$	
5918091809186152461524615231213121312132241322410191101911625716257681768172681726817251423514235142221122112211221122112211221122112211	
59180918091809181524615246152462131221312243122413224191101911071625716257817281728172817514235142351423112021120211202112021120211202112021	0
59180918091826152461524615121312131224132241322411019110191257162571628172817281743514235142352110221102211022110221102211022110221	0
591891801809615245246246153121121321321322422413413221019119110110191625725716716252681781726172682351451423142350211211201101101916257257167162526817817261726823514514231423502112112011011019162572571671625681781726172682351451423142350211211201101101101101101101101101101101	0
591891801809152462461561524213123122121313224224134132219110110191019171625162572571628178172617268514214235235141120220211200000000000000000000000000	0
0 = 9 = 18 = 10 = 10 = 10 = 10 = 10 = 10 = 10	0
59181809918061524261515246312121312244413222413101911101919110625771625877162881726285141235514220211120220000000000000000000000000	0
59 18 18 0 9 9 18 0 15 24 6 6 15 24 24 6 15 21 3 12 21 3 3 12 21 3 22 4 4 13 22 24 13 19 1 10 10 19 1 1 1 10 19 7 16 25 25 7 16 16 25 7 26 8 17 17 26 8 8 17 26 5 14 23 25 14 14 23 5 11 20 2 2 11 20 20 10 10 10 10 10 10 10 10 10 10 10 10 10	0
5918180991802461515246615241213312121312244413222413110191911010191257161625771625268171726881726142355142235142110202111010191257161625771625268171726881726142355142235142110202111010191125716162577162526817172688172614235514223514211020211101019112571616257716252681717268817261423551422351421102021110101911257161625771625268171726881726142355142212110202111010191125716162577162526817172688172614235514221211020211101019112571616257716252681717268817261423551422121102021110101911257161625771625268171726881726142355142212110202111010191125716162577162526817172688172614235514221211020211101019112571616257716252681717268817261423551422121102021110101911257161625771625268171726881726142355142215421102021110019111010191125716162577716526817172688172614235514221511020211100000000000000000000000	0
518901890189624156241532123212321213422134221019101910191012567256725672567256725672567256725672567	0
$\begin{smallmatrix} 18 & 9 & 0 \\ 18 & 0 & 0 \\ 1$	0
5 18 9 0 18 9 0 18 9 24 15 6 24 15 6 12 3 21 12 3 21 12 3 21 3 4 22 13 4 22 1 4 22 1 19 10 1 19 10 1 19 10 21 6 7 25 16 7 26 17 8 26 17 8 26 17 8 14 5 23 14 5 23 14 5 23 2 20 11 2	0
$\begin{array}{c} 18 \\ 9 \\ 9 \\ 0 \\ 18 \\ 18 \\ 9 \\ 0 \\ 6 \\ 24 \\ 15 \\ 6 \\ 24 \\ 15 \\ 6 \\ 3 \\ 21 \\ 12 \\ 12 \\ 3 \\ 21 \\ 12 \\ 13 \\ 14 \\ 22 \\ 21 \\ 14 \\ 4 \\ 22 \\ 10 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 7 \\ 25 \\ 21 \\ 16 \\ 21 \\ 8 \\ 8 \\ 21 \\ 17 \\ 8 \\ 8 \\ 21 \\ 17 \\ 8 \\ 22 \\ 14 \\ 5 \\ 5 \\ 21 \\ 14 \\ 15 \\ 22 \\ 11 \\ 2 \\ 20 \\ 11 \\ 2 \\ 20 \\ 11 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 1 \\ 10 \\ 10$	0
$\begin{smallmatrix} 1899 \\ 01818 \\ 9 \\ 015 \\ 6 \\ 2424 \\ 15 \\ 6 \\ 6 \\ 2415 \\ 121 \\ 2 \\ 3 \\ 2112 \\ 3 \\ 212 \\ 12 \\ 3 \\ 213 \\ 4 \\ 222 \\ 13 \\ 4 \\ 4 \\ 221 \\ 19 \\ 10 \\ 1 \\ 19 \\ 10 \\ 1 \\ 19 \\ 7 \\ 2516 \\ 16 \\ 7 \\ 252 \\ 16 \\ 7 \\ 261 \\ 8 \\ 8 \\ 267 \\ 17 \\ 8 \\ 26 \\ 5 \\ 231 \\ 14 \\ 5 \\ 23 \\ 214 \\ 5 \\ 12 \\ 20 \\ 11 \\ 2 \\ 2 \\ 20 \\ 11 \\ 2 \\ 2 \\ 2 \\ 10 \\ 1 \\ 1 \\ 10 \\ 10$	0
$\begin{array}{c} 18 \\ 9 \\ 9 \\ 0 \\ 18 \\ 18 \\ 9 \\ 0 \\ 24 \\ 15 \\ 6 \\ 24 \\ 15 \\ 6 \\ 24 \\ 12 \\ 3 \\ 21 \\ 21 \\ 23 \\ 21 \\ 21 \\ 14 \\ 22 \\ 21 \\ 14 \\ 4 \\ 22 \\ 13 \\ 1 \\ 19 \\ 10 \\ 1 \\ 19 \\ 10 \\ 1 \\ 19 \\ 10 \\ 1 \\ 216 \\ 7 \\ 7 \\ 256 \\ 17 \\ 8 \\ 8 \\ 267 \\ 18 \\ 8 \\ 267 \\ 17 \\ 8 \\ 264 \\ 15 \\ 23 \\ 214 \\ 5 \\ 5 \\ 23 \\ 14 \\ 2 \\ 20 \\ 111 \\ 2 \\ 20 \\ 10 \\ 1$	0
5 18 9 18 9 0 9 0 18 6 24 5 24 5 6 15 6 24 3 21 21 21 23 21 3 4 22 4 22 13 22 13 4 10 1 19 1 19 10 19 10 1 16 7 25 7 25 16 25 16 7 26 7 8 17 8 26 8 26 17 23 14 5 14 5 23 5 23 14 20 11 2 11 2 20 2 20 20 20 20 20 20 20 20 20 20 20	0
$\begin{smallmatrix} 18918909018156246241524156212312321321213422422132134191011011911910725162167167226178178268261752314231451452311220220122012201220122012201220122012$	0
$\begin{smallmatrix} 18 & 9 & 18 & 9 & 0 & 9 & 0 & 182415 & 6 & 15 & 6 & 24 & 6 & 24 & 15 & 12 & 3 & 3 & 21 & 22 & 12 & 21 & 3 & 13 & 4 & 22 & 4 & 22 & 13 & 22 & 13 & 4 & 1 & 19 & 10 & 10 & 1 & 10 & 1 & 19 & 25 & 67 & 7 & 25 & 7 & 25 & 16 & 26 & 78 & 178 & 26 & 8 & 26 & 71 & 4 & 5 & 23 & 5 & 23 & 14 & 23 & 14 & 5 & 2 & 20 & 11 & 2 & 11 & 21 & 21 & 21 &$	0
59 18 0 9 18 0 9 18 15 24 6 15 24 6 15 24 6 21 3 12 21 3 12 21 3 12 4 13 22 4 13 22 10 19 1 10 19 1 25 7 16 25 7 16 8 17 26 17 26	0
5918918018091524624615615242131231211213413221324224131019119110110192571671625162578172617268268171423523514514230211211201101101925716716251625781726172682681714235235145142302112112011011019257167162516257817261726826817142352351451423021121120110110192571671625162578172617268268171423523514514230211211201101101925716716251625781726172682681714235235145142302112112011011019257167162516257817261726826817142352351451423021121120110110192571671625162578172617268268171423523514514230211211201101101925716716251625781726172682681714235235145142302112112011011019257167162516257817261726826817142352351451423021121120110110192571671625162578172617268268171423523514514230211211201101101925716716251625781726172682681714235235145142302112112011011019257167162516257817261726826817142352351451423021121120110101010000000000000000000	0
59181809918015246615242661521312121331221413222413132241019111019191102571616257716258172626817172681423551423235142021111202200000000000000000000000000	0
5 18 9 0 18 9 0 18 9 15 6 24 15 6 24 21 23 21 23 21 24 22 13 4 22 13 14 22 13 10 1 19 10 1 19 20 17 25 16 7 25 16 7 8 26 77 8 27 77 8 27 77 77 77 77 77 77 77	0
$\begin{array}{c} 18 \\ 9 \\ 9 \\ 0 \\ 18 \\ 18 \\ 9 \\ 0 \\ 15 \\ 6 \\ 24 \\ 215 \\ 6 \\ 6 \\ 24 \\ 15 \\ 12 \\ 12 \\ 3 \\ 3 \\ 21 \\ 12 \\ 3 \\ 21 \\ 12 \\ 3 \\ 21 \\ 4 \\ 22 \\ 13 \\ 14 \\ 22 \\ 21 \\ 3 \\ 14 \\ 22 \\ 21 \\ 3 \\ 14 \\ 22 \\ 21 \\ 3 \\ 14 \\ 22 \\ 21 \\ 3 \\ 14 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 21 \\ 7 \\ 7 \\ 25 \\ 16 \\ 16 \\ 7 \\ 25 \\ 8 \\ 26 \\ 17 \\ 7 \\ 8 \\ 26 \\ 21 \\ 7 \\ 8 \\ 26 \\ 17 \\ 7 \\ 8 \\ 26 \\ 21 \\ 18 \\ 14 \\ 5 \\ 23 \\ 21 \\ 12 \\ 2 \\ 21 \\ 13 \\ 14 \\ 22 \\ 21 \\ 13 \\ 4 \\ 22 \\ 21 \\ 14 \\ 10 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 19 \\ 10 \\ 1 \\ 1 \\ 10 \\ 10$	0
51891890901815624624152415621232132132134134221011911910101251671672572518261726178178261452352314231452011211220200000000000000000000000000	0
59 18 9 18 0 18 0 9 6 15 24 15 24 6 24 6 15 3 12 12 12 13 21 3 21 3 12 4 13 22 13 22 4 22 4	0
5 18 9 9 0 18 18 9 0 6 24 15 15 6 24 24 15 6 3 21 21 21 3 4 22 13 13 4 22 22 13 4 1 19 10 1 1 19 19 10 1 7 25 16 16 7 25 25 16 7 8 26 71 7 8 26 26 7 8 5 23 14 14 5 23 23 14 5 2 20 11 11 2 20 20 20 20 20 20 20 20 20 20 20 20 2	0
59 18 09 18 09 18 6 15 24 6 15 24 3 12 13 12 13 12 14 13 22 4 13 22 4 13 22 1 10 19 1 10 19 7 16 25 7 16 25 8 17 26 8 17 26 8 17 26 5 14 23 5 14 23 5 14 23 2 11 20 2 11 20 2 12 20 2	~

The constructed design D is an SCD(81, $(3^2)^5$, 27), where D₁ is an SOA(81, 5, 3^2 , 2+) and D₂ is an LHD(81, 27). When D₁ is collapsed into three levels, the rows in D₂ corresponding to each level of any factor in the collapsed D₁ form an LHD(27, 27). Furthermore, when D₁ is collapsed into 9 × 3 or 3 × 9 levels, the rows in D₂ corresponding to each level combination of any two factors in the collapsed D₁ form an LHD(3, 27). Moreover, \tilde{d}_1 and \tilde{d}_2 achieve stratification on 3 × 3 grids.

Table 5 presents several SCDs constructed using the method in Algorithm 2. The designs ensure good stratification for quantitative factors while accommodating varying numbers of qualitative and quantitative factors. As the run size of the experimental design increases, the number of columns for qualitative and quantitative factors also increases. This highlights the SCD's effectiveness in achieving space-filling and its capacity to handle multi-factor experiments efficiently.

s	п	D_1	D_2	p+q
2	16	$SOA(16, 2, 2^2, 2+)$	LHD(16,8)	10
2	32	$SOA(32, 2, 2^2, 2+)$	LHD(32, 16)	18
2	64	$SOA(64, 2, 2^2, 2+)$	LHD(64, 24)	26
2	128	$SOA(128, 2, 2^2, 2+)$	LHD(128, 32)	34
3	81	$SOA(81, 5, 3^2, 2+)$	LHD(81,27)	32
3	243	$SOA(243, 5, 3^2, 2+)$	LHD(243,54)	59
3	729	$SOA(729, 5, 3^2, 2+)$	LHD(729,81)	86
4	256	$SOA(256, 7, 4^2, 2+)$	LHD(256,64)	71
4	1024	$SOA(1024, 7, 4^2, 2+)$	LHD(1024, 128)	135
5	625	$SOA(625, 9, 5^2, 2+)$	LHD(625,125)	134

Table 5. Some SCDs generated by Algorithm 2.

4. Concluding Remarks

In this paper, a new class of designs for computer experiments with both quantitative and qualitative factors, referred to as SCDs, is introduced. These designs can be seen as a generalization of MCDs proposed by Deng, Hung and Lin (2015) [20] and DCDs introduced by Yang et al. (2023) [25]. Similar to DCDs, certain stratification can be achieved between any two qualitative factors and all quantitative factors in SCDs. The stratification requirements among qualitative factor columns are relaxed, making the run sizes of SCDs more flexible than those of DCDs.

The characteristics and construction methods of SCDs are investigated. Necessary and sufficient conditions for the existence of SCDs are derived from three perspectives, providing valuable insights for their construction. Additionally, two methods for constructing SCDs are proposed. When $s \ge 2$ is a prime or a prime power, the constructed SCDs can accommodate 2s - 1 qualitative factors and a substantial number of quantitative factors. Furthermore, a series of SCDs with s^u rows and $(u - 3)s^3$ columns of quantitative factors, where $u \ge 4$, is constructed, with certain columns of quantitative factors achieving stratification in two or higher dimensions. In the future, designs for qualitative factors could be extended to achieve higher-dimensional stratification properties, such as SOAs with strength $t \ge 3$ or SOAs with strength 3-.

Author Contributions: M.-M.L.: conceptualization, methodology, formal analysis, writing—original draft preparation; M.-Q.L.: methodology, supervision, writing—review and editing, funding acquisition; J.-Y.Y.: methodology, formal analysis, writing—review and editing, supervision. All authors have read and agreed to the published version of the manuscript.

Funding: This work was supported by the National Natural Science Foundation of China (Grant Nos. 12131001 and 12371260).

13 of 15

Data Availability Statement: The original contributions presented in the study are included in the article; further inquiries can be directed to the corresponding author.

Conflicts of Interest: The authors declare no conflicts of interest.

Appendix A

Proof of Theorem 1. According to part (ii) of Definition 1, the n/s rows in D_2 corresponding to each level of any factor in \overline{D}_1 form an LHD if and only if any two columns in \overline{D}_1 and \overline{D}_2 together form an OA of strength 2. In other words, for any $1 \le i \le q, 1 \le k \le p$, the pair $(\overline{z}_i, \overline{d}_k)$ forms an OA(n, 2, s(n/s), 2). Similarly, when any two columns of D_1 are collapsed into $s^2 \times s$ or $s \times s^2$ levels, the n/s^3 rows in D_2 corresponding to each level combination of any two factors in the collapsed D_1 form an LHD if and only if $(z_i, \overline{z}_j, \overline{d}_k)$ form an OA of strength 3, where $1 \le i \ne j \le q$ and $1 \le k \le p$. This completes the proof. \Box

Proof of Theorem 2. The necessity is first shown. By Theorem 1, an SCD (n, s^q, p) exists if and only if both conditions (i) and (ii) of Theorem 1 are satisfied. The fact that $(z_i = s\bar{z}_i + b_i, \bar{z}_j, \tilde{d}_k)$ forms an OA $(n, 3, s^2s(n/s^3), 3)$ implies that $(\bar{z}_i, b_i, \bar{z}_j, \tilde{d}_k)$ forms an OA $(n, 3, s^3(n/s^3), 4)$. Consequently, $(\bar{z}_i, b_i, \bar{z}_j)$ can be divided into n/s^3 full factorials, denoted by $(\bar{z}_i^{(l)}, b_i^{(l)}, \bar{z}_j^{(l)})$, for $l = 1, ..., n/s^3$, with each corresponding to one level in \tilde{d}_k . Since (\bar{z}_i, \bar{d}_k) is an OA(n, 2, s(n/s), 2), we have $\tilde{d}_k = \lfloor \bar{d}_k/s^2 \rfloor$, and thus each $A^{(l)} = (\bar{z}_1^{(l)}, ..., \bar{z}_q^{(l)})$ is a CROA $(s^3, q, s, 2)$.

Now, we prove the sufficiency by providing a method to construct D_2 when $D_1 = sA + B$, with A and B satisfying the two conditions in Theorem 2. According to Theorem 2, $A = (A^{(1)^T}, \ldots, A^{(n/s^3)^T})^T$, where $A^{(l)} = (A_1^{(l)}, \ldots, A_{s^2}^{(l)})$, is a CROA $(s^3, q, s, 2)$ for $1 \leq l \leq n/s^3$. Define $\overline{D}_2 = (s^2c_1 + E, \ldots, s^2c_{(n/s^3)!} + E)$, where $c_i = \omega_i \otimes 1_{s^3}$, with ω_i being a permutation of $(0, \ldots, n/s^3 - 1)$, for $1 \leq i \leq (n/s^3)!$. $E = (e_1, \ldots, e_{(s^{2!})^{n/s^3}})$, where $e_k = (e_k^{(1)^T}, \ldots, e_k^{(n/s^3)^T})^T$, and $e_k^{(l)} = ((\beta_{1,k}^{(l)} \otimes 1_s)^T, \ldots, (\beta_{s,k}^{(l)} \otimes 1_s)^T)^T$, with $(\beta_{1,k}^{(l)^T}, \ldots, \beta_{s,k}^{(l)^T})^T$ being a permutation of $(0, \ldots, s^2 - 1)^T$, for $1 \leq k \leq (s^2)!^{n/s^3}$. Clearly, (\bar{z}_i, \bar{d}_k) forms an OA of strength t = 2. Since $\tilde{d}_k = (c_1, \ldots, c_{(n/s^3)!}), (\bar{z}_i^{(l)}, \bar{z}_j^{(l)}, b_i^{(l)})$ forms an OA $(n, 3, s^3, 3)$, with each corresponding to one level in \tilde{d}_k . Therefore, $(\bar{z}_i, \bar{z}_j, \tilde{d}_k)$ forms an OA of strength 3. By Theorem 1, an SCD exists. The proof is completed. \Box

Proof of Theorem 3. By Theorem 1, an SCD (n, s^q, p) exists if and only if both conditions (i) and (ii) in Theorem 1 are satisfied. Since $\overline{D}_2 = s^2C + sE + F$, for any $1 \le i \le q$ and $1 \le k \le p$, $(\overline{z}_i, \overline{d}_k)$ forms an OA(n, 2, s(n/s), 2) if and only if $(\overline{z}_i, f_k, e_k, \overline{d}_k)$ forms an OA $(n, 4, s^3(n/s^3), 4)$. Similarly, for $D_1 = sA + B$, $(z_i, \overline{z}_j, \overline{d}_k)$ forms an OA $(n, 3, s^2s(n/s^3), 3)$ if and only if $(\overline{z}_i, b_i, \overline{z}_i, \overline{d}_k)$ forms an OA $(n, 4, s^3(n/s^3), 4)$. The proof is completed. \Box

Proof of Theorem 4. Let $O = (0, ..., s - 1)^T \oplus H = (o_1, ..., o_s)$ be a CROA(n, s, s, 2), where $H = (h_1, ..., h_s)$ is a D(s, s, s). Define $F_i = H \oplus o_i$ for i = 1, ..., s, and let $F = (F_1, ..., F_s)$. The following two lemmas are first presented.

Lemma A1. If *O* is a CROA(n, q, s, 2) and *H* is a D(s, s, s), then $O \oplus H$ is a CROA(ns, qs, s, 2).

Lemma A2. Define a_1 as the first column of F_1 , and A_i as the first and *i*-th columns of F_i for i = 1, ..., s. Then, $A = (a_1, A_1, ..., A_s) = (a_1, ..., a_{2s-1})$. For $1 \leq j_{i,m}$, $k_{i,m} \leq s$, where $1 \leq i \leq s$, $1 \leq m \leq 2$, $j_{i,m} \neq k_{i,m}$, and $j_{i,m} \neq 1$, define $b_1 = h_{j_{1,1}} \oplus o_{k_{1,1}}$ with $j_{1,1} \neq 1$, $b_{i1} = h_{j_{i,1}} \oplus o_{k_{i,1}}$ with $k_{i,1} \neq 1$ and $j_{i,1} \neq i$, or $k_{i,1} = 1$ and $j_{i,1} = i$, and $b_{i2} = h_{j_{i,2}} \oplus o_{k_{i,2}}$ with $j_{i,2} = i$. Let $B_i = (b_{i1}, b_{i2})$ and $B = (b_1, B_2, ..., B_s) = (b_1, ..., b_{2s-1})$. Then, (a_i, a_j, b_i) forms an OA(n, 3, s, 3) for any $1 \leq i \neq j \leq 2s - 1$.

Based on the proof of Theorem 2, combined with Lemma A1 and Lemma A2, the design $D = (D_1, D_2)$ constructed in Algorithm 1 is an SCD. \Box

Proof of Theorem 5. We first present Lemma A3, which is a straightforward generalization of Theorem 3.29 from Hedayat, Sloane and Stufken (1999) [27].

Lemma A3. For A, B, C, E, F, and R_1, \ldots, R_{u-3} in Algorithm 2, we have that

- (i) $(\bar{z}_i, f_k, c_k, r_{1,f}, \dots, r_{u-3,f})$ is an $OA(s^u, u, s, u)$, for any $i \neq k, 1 \leq f \leq s^3$;
- (ii) $(\bar{z}_i, e_i, \bar{z}_j, r_{1,f}, \dots, r_{u-3,f})$ is an $OA(s^u, u, s, u)$, for any $i \neq j, 1 \leq f \leq s^3$;
- (iii) $(r_{1,f}, \ldots, r_{u-3,f})$ is an $OA(s^u, u-3, s, u-3)$;
- (iv) $(r_{1,f}, ..., r_{u-3,f}, r_{v,l})$ is an $OA(s^u, u-2, s, u-2)$ for any $1 \le v \le u-3, 1 \le f \ne l \le s^3$.

From the proof of Theorem 4 in He, Cheng and Tang (2018), D_1 is an SOA. From Theorem 3 and conditions (i) and (ii) of Lemma A3, we know that the constructed design D is an SCD. The number of quantitative factors that the SCD can accommodate is equal to the number of columns in B, so $p = (u - 3)s^3$. The number of qualitative factors is equal to the number of columns in S_1 minus one, which leads to the conclusion of (i) in Theorem 5. For $\lfloor (i_1 - 1)/(u - 3) \rfloor = \cdots = \lfloor (i_m - 1)/(u - 3) \rfloor$, where $m \leq u - 3$, the columns $\tilde{d}_{i_1}, \ldots, \tilde{d}_{i_m}$ come from the same B_f , with $1 \leq f \leq s^3$. From (iii) of Lemma A3, condition (ii) of Theorem 5 is established. When $\lfloor (i - 1)/(u - 3) \rfloor \neq \lfloor (i' - 1)/(u - 3) \rfloor$, \tilde{d}_i and $\tilde{d}_{i'}$ are columns from B_u and B_v , respectively, with $u \neq v$. From (iv) of Lemma A3, condition (iii) of Theorem 5 is established. \Box

References

- 1. Santner, T.J.; Williams, B.J.; Notz, W.I. The Design and Analysis of Computer Experiments; Springer: New York, NY, USA, 2003.
- 2. Fang, K.T.; Li, R.; Sudjianto, A. Design and Modeling for Computer Experiments; Chapman & Hall/CRC: New York, NY, USA, 2006.
- 3. Wu, C.F.J.; Hamada, M.S. Experiments: Planning, Analysis and Optimization, 3rd ed.; John Wiley & Sons: New York, NY, USA, 2021.
- McKay, M.D.; Beckman, R.J.; Conover, W.J. A comparison of three methods for selecting values of input variables in the analysis of output from a computer code. *Technometrics* 1979, 21, 239–245.
- 5. Owen, A.B. Orthogonal arrays for computer experiments, integration, and visualization. *Stat. Sin.* **1992**, *2*, 439–452.
- 6. Tang, B. Orthogonal array-based Latin hypercubes. J. Am. Stat. Assoc. 1993, 88, 1392–1397. [CrossRef]
- He, Y.; Tang, B. Strong orthogonal arrays and associated Latin hypercubes for computer experiments. *Biometrika* 2013, 100, 254–260. [CrossRef]
- 8. He, Y.; Tang, B. A characterization of strong orthogonal arrays of strength three. Ann. Stat. 2014, 42, 1347–1360. [CrossRef]
- 9. He, Y.; Cheng, C.S.; Tang, B. Strong orthogonal arrays of strength two plus. *Ann. Stat.* **2018**, *46*, 457–468. [CrossRef]
- Zhou, Y.; Tang, B. Column-orthogonal strong orthogonal arrays of strength two plus and three minus. *Biometrika* 2019, 106, 997–1004. [CrossRef]
- 11. Liu, H.; Sun, F.; Lin, D.K.J.; Liu, M.Q. On construction of mappable nearly orthogonal arrays with column orthogonality. *Commun. Math. Stat.* **2024**, *in press*. [CrossRef]
- 12. Wang, C.; Lin, D.K.J.; Liu, M.Q. A new class of orthogonal designs with good low dimensional space-filling properties. *Stat. Sin.* **2024**, *in press*. [CrossRef]
- 13. Qian, P.Z.G.; Wu, H.; Wu, C.F.J. Gaussian process models for computer experiments with qualitative and quantitative factors. *Technometrics* **2008**, *50*, 383–396. [CrossRef]
- 14. Han, G.; Santner, T.J.; Notz, W.I.; Bartel, D. Prediction for computer experiments having quantitative and qualitative input variables. *Technometrics* **2009**, *51*, 278–288. [CrossRef]
- 15. Huang, H.Z.; Lin, D.K.J.; Liu, M.Q.; Yang, J.F. Computer experiments with both qualitative and quantitative variables. *Technometrics* **2016**, *58*, 495–507. [CrossRef]
- 16. Qian, P.Z.G. Sliced Latin hypercube designs. J. Am. Stat. Assoc. 2012, 107, 393–399. [CrossRef]
- 17. Kumar, A.A.; Mandal, B.N.; Parsad, R.; Dash, S.; Kumar, M. On construction of sliced orthogonal Latin hypercube designs. *J. Stat. Theory Pract.* **2024**, *in press*. [CrossRef]
- Guo, B.; Li, X.; Liu, M.Q.; Yang, X. Construction of orthogonal general sliced Latin hypercube designs. *Stat Pap.* 2023, 64, 987–1014. [CrossRef]

- Wang, Y.; Wang, D.; Yue, X. Sequential Latin hypercube design for two-layer computer simulators. J. Qual. Technol. 2023, 56, 71–85. [CrossRef]
- 20. Deng, X.; Hung, Y.; Lin, C.D. Design for computer experiments with qualitative and quantitative factors. *Stat. Sin.* **2015**, *25*, 1567–1581. [CrossRef]
- 21. He, Y.; Lin, C.D., Sun, F.; Lv, B.J. Marginally coupled designs for two-level qualitative factors. *J. Stat. Plan. Inference* **2017**, *187*, 103–108. [CrossRef]
- 22. He, Y.; Lin, C.D.; Sun, F. On the construction of marginally coupled designs. Stat. Sin. 2017, 27, 665–683. [CrossRef]
- 23. He, Y.; Lin, C.D.; Sun, F. Construction of marginally coupled designs by subspace theory. Bernoulli 2019, 25, 2163–2182. [CrossRef]
- 24. Zhou, W.; Yang, J.; Liu, M.Q. Construction of orthogonal marginally coupled designs. Stat. Pap. 2021, 62, 1795–1820. [CrossRef]
- 25. Yang, F.; Lin, C.D.; Zhou, Y.; He, Y. Doubly coupled designs for computer experiments with both qualitative and quantitative factors. *Stat. Sin.* **2023**, *33*, 1923–1942. [CrossRef]
- 26. Zhou, W.; Huang, S.; Li, M. Group doubly coupled designs. *Mathematics* **2024**, 12, 1352. [CrossRef]
- 27. Hedayat, A.S.; Sloane, N.J.A.; Stufken, J. Orthogonal Arrays: Theory and Applications; Springer: New York, NY, USA, 1999.

Disclaimer/Publisher's Note: The statements, opinions and data contained in all publications are solely those of the individual author(s) and contributor(s) and not of MDPI and/or the editor(s). MDPI and/or the editor(s) disclaim responsibility for any injury to people or property resulting from any ideas, methods, instructions or products referred to in the content.